# Approximability, compactness and random dense sequences

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### **Notation**

X separable metric space

$$\mathcal{R}_X := \{ \rho \, | \, \rho : \subseteq \mathbb{N}^{\,\mathbb{N}} \to X \text{ onto} \}$$

$$\mathcal{I}_{\mathsf{d}} := \{ \nu \mid \nu : \subseteq \mathbb{N} \to X \text{ dense} \}$$

Fix open cover  $(U_i)_1^r$ ,  $[r] := \{1, ..., r\}$ 

$$R:X \Longrightarrow [r], x \mapsto \{i \mid U_i \ni x\}$$

$$S:[r] \Rightarrow X, i \mapsto U_i$$

$$T_{\epsilon}: X \rightrightarrows X, x \mapsto \{y \mid d(x,y) < \epsilon\} \ (\epsilon > 0)$$

$$\delta \leq_{\mathsf{a}} \rho : \iff (\forall \epsilon > 0) (T_{\epsilon} \text{ is } (\delta, \rho)\text{-computable })$$
  
 $\delta \leq \rho \iff (\mathsf{id}_X : X \to X \text{ is } (\delta, \rho)\text{-computable })$ 

 $\delta \in \mathcal{R}_X$  compact if for each finite open cover  $(U_i)_1^r$ ,  $R:X \rightrightarrows [r]$  is  $(\delta, \delta_{\mathbb{N}}|^{[r]})$ -computable

**Example:** If X compact and  $\nu \in \mathcal{I}_d$ , the *stan-dard representation* (equivalent to Cauchy representation in case of a computable metric space) defined by

$$p \in \delta^{-1}\{x\} : \iff \{p_i - 1 \mid i \in \mathbb{N} \land p_i \ge 1\} = \{k \mid \alpha(k) \ni x\},$$
  
$$\alpha : \mathbb{N} \to \mathcal{T}_X, \langle i, j \rangle \mapsto B_d(\nu(i); \nu_{\mathbb{O}} + (j)).$$

**Lemma 1.** Suppose X compact,  $\delta, \rho \in \mathcal{R}_X$ ,  $\delta$  compact. If  $\rho$ -computable points are dense in X then  $\delta \leq_a \rho$ .

**Proof.** Given  $\epsilon > 0$ , pick open cover  $(U_i)_1^r$  with  $\max_i \operatorname{diam} U_i < \epsilon$ . Then  $(\delta, \delta_{\mathbb{N}}|^{[r]})$ -computability of R and  $(\delta_{\mathbb{N}}|^{[r]}, \rho)$ -computability of S implies  $(\delta, \rho)$ -computability of  $T_{\epsilon}$ .

Some generalisation is possible. Consider e.g. **Lemma 2.** If (X,d) a totally bounded metric space and  $\nu \in \mathcal{I}_d$  then for any  $r \in \mathbb{Q}^+$  there exists finite  $A \subseteq \text{dom } \nu$  with  $X = \bigcup_{a \in A} B_d(\nu(a); r)$ .

For  $\nu_0, \nu_1, \nu, \lambda \in \mathcal{I}_d$ , write

$$\nu \leq_{\mathsf{a}} \lambda : \iff (\forall \epsilon > 0)(\exists h \in P^{(1)})(\forall c) \Big( c \in \mathsf{dom}\, \nu \implies c \in (\lambda \circ h)^{-1} B(\nu(c); \epsilon) \Big).$$
$$\mathsf{dom}(\nu_0 \oplus \nu_1) := \dot{\cup}_i (2\,\mathsf{dom}\, \nu_i + i), \quad \nu_0 \oplus \nu_1(2a+i) := \nu_i(a)$$

 $\nu \in \mathcal{I}_{d}$  compact if any finite open cover  $(U_{i})_{1}^{r}$  admits some  $f \in P^{(1)}$  with

$$\operatorname{dom} \nu \subseteq f^{-1}[r] \wedge (\forall a \in \operatorname{dom} \nu) \left( \nu(a) \in U_{f(a)} \right) \tag{1}$$

**Proposition 3.** Let X be a separable metric space,  $\rho_{\nu}$  the Cauchy representation for  $\nu \in \mathcal{I}_d$ . For any  $\nu, \lambda \in \mathcal{I}_d$  and  $\delta, \rho \in \mathcal{R}_X$ ,

- 1.  $\nu \leq_a \lambda \iff \rho_{\nu} \leq_a \rho_{\lambda}$
- 2.  $\delta \leq \rho \implies \delta \leq_{a} \rho$
- 3.  $\delta \sqcup \rho$  is a least upper bound of  $\{\delta, \rho\}$  w.r.t.  $\leq_a$
- 4.  $\nu \oplus \lambda$  is a least upper bound of  $\{\nu, \lambda\}$  w.r.t.  $\leq_a$

**Proof of (3):** First apply (2) in  $\delta_i \leq \delta_0 \sqcup \delta_1$  (i < 2). If also  $\delta_i \leq_a \rho$ , say via  $F_i$  at precision  $\epsilon$  (i < 2), then  $\delta_0 \sqcup \delta_1 \leq_a \rho$  at precision  $\epsilon$  via  $F : \subseteq \mathbb{N}^{\mathbb{N}} \to \mathbb{N}^{\mathbb{N}}, i.p \mapsto F_i(p)$ .

Each  $\leq_a$  also reflexive and transitive. Thus above operations give rise to upper semilattice structures, on  $\mathcal{R}_X/\equiv_a$ ,  $I_a:=\mathcal{I}_d/\equiv_a$  and  $(\mathcal{I}_d\cap X^\mathbb{N})/\equiv_a$ . If X compact, analogue of Lemma 1 for dense partial sequences implies compact  $\nu\in\mathcal{I}_d$  form a least element of  $I_a$  (later we show there exists compact  $\nu\in\mathcal{I}_d$ ).

**Proposition 4.** 1.  $\rho \leq \angle a \rho$ 

- 2.  $\rho \leq_a \rho_{Cf}$
- 3.  $\rho_{<} \leq_a \rho_{<}$  and symmetrically (replace  $<, \leq by >, \geq$ ).

Hence  $\rho \leq_a \rho_{Cf}, \rho \leq_. \rho_{\geq}$  and  $\rho_{b,n} \leq_a \rho_{Cf}, \rho \geq_. \rho_{\leq}$  and (3) are the only  $\leq_a$ -reductions not shown in the left figure.

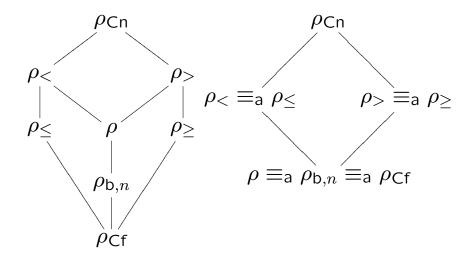
**Proof.** (1): Let F realise  $\rho_{\leq} \leq_{\mathsf{a}} \rho_{>}$  to precision  $\epsilon$ ; where  $p \in \rho_{\leq}^{-1}\{x\}$  enumerates all rationals  $\leq x$ , q = F(p) enumerates strict right Dedekind cut of some y s.t.  $|x-y| < \epsilon$ .  $q_0$  is output after finitely many steps, with only finite prefix  $p^N$  of p read from input. Let  $p' \in \rho_{\leq}^{-1}\{z\}$  for some  $z \geq \nu_{\mathbb{Q}}(q_0) + \epsilon$  (>  $y + \epsilon > x$ ) with  $(p')^N = p^N$ . Then  $F(p')_0 = q_0$ , so  $(\rho_{<} \circ F)(p') < \nu_{\mathbb{Q}}(q_0) \leq z - \epsilon$ , contradiction.

(3): Recall the  $\mathbf{T}_0$ -topology  $\tau_< = \{(x, \infty) \mid x \in \mathbb{R}\} \cup \{\emptyset, \mathbb{R}\}$  on  $\mathbb{R}$  (with respect to which  $\rho_<$  is admissible) and the following

**Lemma 5.** For any  $D \subseteq \mathbb{R}$ , a function  $f :\subseteq \mathbb{R} \to \mathbb{R}$  with dom f = D is  $(\tau_{<}, \tau_{<})$ -continuous iff it is left-continuous and nondecreasing.

We specify f which is  $(\rho_<, \rho_\le)$ -computable and lies in the open  $\epsilon$ -envelope of  $\mathrm{id}_\mathbb{R}$ . First, consider  $f:\mathbb{R}\to\mathbb{R}$  as in Lemma 5 which is also piecewise constant: take  $f(t)=c_i$  for  $t_i< t\le t_{i+1}$  where strictly increasing  $(c_i)_{i\in\mathbb{Z}}, (t_i)_{i\in\mathbb{Z}}$  have  $\inf t_i=-\infty$ ,  $\sup t_i=\infty$ ,  $c_i-\epsilon\le t_i\wedge t_{i+1}< c_i+\epsilon$   $(i\in\mathbb{Z}).$  If  $(t_i)_i$  are uniformly right-computable and  $(c_i)_i$  uniformly left-computable then f is  $(\rho_<,\rho_<)$ -computable. If  $(c_i)_{i\in\mathbb{Z}}\subseteq\mathbb{R}\setminus\mathbb{Q}$  then any  $(\rho_<,\rho_<)$ -realiser of f also  $(\rho_<,\rho_\le)$ -realises it. For instance, if  $\epsilon=2^{-j+\frac{1}{2}}$  we can take  $t_n:=n\epsilon$ ,  $c_n:=t_n+\frac{\epsilon}{2}=\frac{2n+1}{2}\epsilon$   $(\not\in\mathbb{Q}),\ n\in\mathbb{Z}$ .

Note now that (1) (with transitivity of  $\leq_a$ ) implies  $\delta_0 \not\leq_a$   $\delta_1$  for all  $(\delta_0, \delta_1) \in \{\rho_\leq, \rho_<, \rho_{\mathsf{Cn}}\} \times \{\rho_{\mathsf{Cf}}, \rho_{\mathsf{b},n}, \rho_\geq, \rho, \rho_>\}$ . With symmetric version (exchanging  $<, \leq$  with  $>, \geq$ ), shows  $\rho_{\mathsf{Cn}} \not\leq_a \delta$  for all  $\delta \in S \setminus \{\rho_{\mathsf{Cn}}\}$ , where  $S := \{\rho_{\mathsf{Cf}}, \rho_\leq, \rho_\geq, \rho_{\mathsf{b},n}, \rho, \rho_<, \rho_>, \rho_{\mathsf{Cn}}\}$  (for convenience we fix n), while  $\delta \leq \rho_{\mathsf{Cn}}$  for all  $\delta \in S$ . For  $\delta_0 \in \{\rho_\leq, \rho_\geq\}$  again (1) and figure completely determine either  $\delta_0 \not\leq_a \delta_1$  or  $\delta_0 \leq \delta_1$  as  $\delta_1 \in S$  varies. Case  $\delta_0 \in \{\rho_<, \rho_>\}$  is similar except where (3) and its symmetric version apply. So assume  $\delta_0 \in \{\rho_{\mathsf{Cf}}, \rho_{\mathsf{b},n}, \rho\}$ , and  $\delta_1 \in S \setminus \{\rho_<, \rho_>, \rho_{\mathsf{Cn}}\}$  (otherwise  $\delta_0 \leq \delta_1$ ). By (2) we already have  $\delta_0 \leq_a \delta_1$ , which completes the proof.



From the figure  $(\rho \equiv_a \rho_{b,n} \equiv_a \rho_{Cf})$  and known facts it follows  $\leq_a$  does not imply continuous reducibility  $\leq_t$  (nor does  $\equiv_a$  guarantee the same final topology). The converse inclusion  $(\leq_t) \subseteq (\leq_a)$  also fails in general, using next construction from [?] (with Proposition 3(1)).

**Definition 6.** Let  $(Y, \nu)$  be a numbered set with  $|Y| \ge 2$  and fix  $x, y \in Y$  with  $x \ne y$ . For each  $A \subseteq \mathbb{N}$  write  $\nu^A(2k+i) = \nu^A_{x,y}(2k+i) := \begin{cases} \nu(k), & \text{if } i = 0 \land k \in \text{dom } \nu, \\ x, & \text{if } i = 1 \land k \in A, \\ y, & \text{if } i = 1 \land k \in \mathbb{N} \setminus A \end{cases}$  (i < 2).

Iterating, we can compare towers constructed this way. First consider a generalisation of compact  $\nu \in \mathcal{I}_d$ .

**Definition 7.** Let X be a separable metric space.  $\nu \in \mathcal{I}_d$  separating at finite precision if for any distinct  $x, y \in X$  there exist open cover  $(U_i)_1^r$ ,  $V, W \in \mathcal{T}_X$ ,  $h \in P^{(1)}$  with

 $x \in V \land y \in W \land (\forall c \in \nu^{-1}(V \cup W))(c \in h^{-1}[r] \land \nu(c) \in U_{h(c)}) \land \{i \mid U_i \cap V \neq \emptyset \neq U_i \cap W\} = \emptyset.$ 

**Theorem 8.** Let X be a separable metric space,  $\nu_0, \ldots, \nu_n, \lambda, \lambda' \in \mathcal{I}_d$ ,  $A_i, B \subseteq \mathbb{N}$   $(i < n \in \mathbb{N})$ ,  $E_n := \{x_i, y_i \mid i < n\} \subseteq X$  and  $x, y \in X$  distinct. Suppose  $\nu_0$  separating a.f.p.,  $\nu_{i+1} = (\nu_i)_{x_i,y_i}^{A_i}$  for all i < n and  $\lambda' = \lambda_{x,y}^{B}$ . If  $E_n \cap \{x,y\} = \emptyset$  and B is nonrecursive then  $\lambda' \not\leq_a \nu_n$ .

**Proof.** Fix  $(U_i)_1^r, V, W, h$  as in definition of separation a.f.p., and pick  $\epsilon > 0$  suff. small that  $B(x;\epsilon) \subseteq V \land B(y;\epsilon) \subseteq W \land N_\epsilon(E_n) \cap \{x,y\} = \emptyset$ . Also suppose  $\lambda'|_{2\mathbb{N}+1} \leq_{\mathbf{a}} \nu_n$  at precision  $\epsilon$  via  $f \in P^{(1)}$ . Write  $k \in \mathbb{N}$ ,  $a = f(2k+1) = \sum_{i=0}^n a_i 2^i$  where  $a_n \in \mathbb{N}$ ,  $(a_i)_{i < n} \subseteq \{0,1\}$ . The last requirement on  $\epsilon$  means  $\lambda'(2k+1) \in \{x,y\} \implies \nu_n(a) \in \operatorname{im} \nu_0 \setminus E_n$  and  $a_i = 0$  for i < n (inductively for  $m = n, \ldots, 1$ , use

$$\operatorname{im} \nu_0 \backslash E_n \ni \nu_m (\sum_{j=0}^m b_j 2^j) = \nu_{m-1}^{A_{m-1}} (\sum_{j=1}^m b_j 2^j) = \nu_{m-1} (\sum_{j < m} b_{j+1} 2^j)$$

where  $(b_j)_0^m = (a_i)_0^n, (a_{i+1})_0^{n-1}, \dots, (a_{i+n-1})_0^1)$ . So, we get  $\nu_n(a) = \nu_n(a_n 2^n) = \dots = \nu_1(a_n 2^1) = \nu_0(a_n)$  while  $g: k \mapsto a_n = \left\lfloor 2^{-n} f(2k+1) \right\rfloor$  is computable with im  $g \subseteq \text{dom } \nu_0$ .

 $(\lambda'(2k+1) = x \implies g(k) \in h^{-1}B_0) \wedge (\lambda'(2k+1) = y \implies g(k) \in h^{-1}B_1)$  for all  $k \in \mathbb{N}$  where  $B_0 := \{i \mid B(x;\epsilon) \cap U_i \neq \emptyset\}$  and  $B_1 := \{i \mid B(y;\epsilon) \cap U_i \neq \emptyset\}$ . In particular,  $B \leq_{\mathsf{m}} B_0$  via  $h \circ g$ , which implies B recursive, a contradiction. So,  $\lambda' \not\leq_{\mathsf{a}} \nu_n$ .  $\square$  **Lemma 9.** 1.  $A \leq_{\mathsf{m}} B \implies \nu^A \leq \nu^B$ ,

2. If  $x \neq y \land \{x_i, y_i \mid i < n\} \cap \{x, y\} = \emptyset \land \emptyset \neq A \neq \mathbb{N} \land (\forall i < n)(\nu_{i+1} = (\nu_i)_{x_i, y_i}^{A_i})$  and  $\lambda_{x, y}^B \leq_a (\nu_n)_{x, y}^A$  where  $\nu_0$  separating a.f.p. then  $B \leq_m A$ .

**Proof.** (1): Fix  $f \in R^{(1)}$  such that  $A = f^{-1}B$  and let  $g: \mathbb{N} \to \mathbb{N}$ ,  $2k+i \mapsto \begin{cases} 2k, & \text{if } i = 0, \\ 2f(k) + 1, & \text{if } i = 1 \end{cases}$ . Then one checks  $\nu^A \leq \nu^B$  via g.

(2): Fix  $(U_i)_1^r, V, W, h$  as in definition of separation a.f.p. (for  $x \neq y$ ),  $\epsilon > 0$  such that  $B(x; \epsilon) \subseteq A$ 

 $V \wedge B(y;\epsilon) \subseteq W \wedge N_{\epsilon}(E_n) \cap \{x,y\} = \emptyset \text{ where } E_n := \{x_i,y_i \ | \ i < n\}. \text{ Let } g \in P^{(1)} \text{ witness } \lambda^B \leq_{\mathsf{a}} (\nu_n)^A \text{ at precision } \epsilon \text{ and denote } l : \mathbb{N} \to \mathbb{N}, k \mapsto \left\lfloor \frac{1}{2}g(2k+1) \right\rfloor,$   $C_i := \{k \in \mathbb{N} \ | \ g(2k+1) \equiv i \pmod{2}\} \ (i=0,1). \text{ Also let } B_0 := \{i \ | \ B(x;\epsilon) \cap U_i \neq \emptyset\}, \ B_1 := \{i \ | \ B(y;\epsilon) \cap U_i \neq \emptyset\}, \ \text{and choose } c \in A, \ d \in \mathbb{N} \setminus A, \ m : k \mapsto \left\lfloor 2^{-n}l(k) \right\rfloor$  and  $f : \mathbb{N} \to \mathbb{N}, k \mapsto \begin{cases} c, & \text{if } k \in C_0 \wedge m(k) \in h^{-1}B_0, \\ d, & \text{if } k \in C_0 \wedge m(k) \in h^{-1}B_1, \\ l(k), & \text{if } k \in C_1 \end{cases}$ 

plainly  $C_0, C_1$  are disjoint recursive sets (with union  $\mathbb{N}$ ), so  $f \in P^{(1)}$ . We show  $f \in R^{(1)}$  with  $(\forall k)(k \in B \iff f(k) \in A)$ .

Firstly, if  $k \in C_0$  then

$$\epsilon > d(\lambda^B(2k+1), (\nu_n \circ l)(k))$$

$$\implies (\nu_n \circ l)(k) \in \operatorname{im} \nu_0 \setminus E_n$$

$$\implies l(k) = a_n 2^n$$

where  $a_n = m(k)$ . So  $(\nu_n \circ l)(k) = \nu_n(a_n.2^n) = \cdots = \nu_0(a_n.2^0) = (\nu_0 \circ m)(k)$  (this also shows  $C_0 \subseteq m^{-1} \operatorname{dom} \nu_0$ ). Now definition of  $B_0, B_1$ 

implies  $(\lambda^B(2k+1) = x \implies m(k) \in h^{-1}B_0) \wedge (\lambda^B(2k+1) = y \implies m(k) \in h^{-1}B_1)$ , so  $k \in C_0$  implies  $k \in B \iff f(k) \in A$ . For  $k \in C_1$ , instead  $\lambda^B_{x,y}(2k+1) = ((\nu_n)^A_{x,y} \circ g)(2k+1)$  with  $k \in B \iff l(k) \in A$  (this uses  $d(x,y) \geq \epsilon$ ). So f has the properties required.

In particular, for any  $\nu_n, x, y$  as above, the map  $\alpha: A \mapsto (\nu_n)_{x,y}^A$  induces an embedding of  $\leq_{\mathsf{m}}$ -degrees in  $\leq_{\mathsf{a}}$ -degrees with least element  $[\nu_n]_{\equiv_{\mathsf{a}}}$ .

# Randomness and density

**Definition 10.** Consider a bounded effective metric space  $(X,d,\nu_0)$ . A **denseness test** (A,a,j) has  $A\subseteq \mathbb{N}$  infinite c.e.,  $a\in \text{dom }\nu_0,\ j\in \mathbb{N}$ ;  $\xi\in X^\mathbb{N}$  fails (A,a,j) if  $\xi\in \cap_{l\in A}X^\mathbb{N}\setminus \sigma^{-l}\left(\alpha\langle a,j\rangle\times X^\mathbb{N}\right)$ , passes (A,a,j) if  $\xi\in \cup_{l\in A}\sigma^{-l}\left(\alpha\langle a,j\rangle\times X^\mathbb{N}\right)$ ; these define resp. closed, open sets in  $X^\mathbb{N}$ .

 $\mathcal{I}_{d}^{\mathsf{rand}} := \{ \xi \in X^{\mathbb{N}} \mid \xi \text{ passes all denseness tests} \}.$ 

We introduce another generalisation of the definition of compact sequences.

**Definition 11.** Suppose X a separable metric space,  $(U_i)_1^r$  a finite open cover,  $\epsilon > 0$  with each  $N_{\epsilon}(U_i)$  nondense and  $\nu \in \mathcal{I}_d$ ,  $f \in P^{(1)}$  s.t. (1) holds. Then  $\nu$  is properly covering.

If  $\nu \in \mathcal{I}_d$  is properly covering then  $|X| \geq 2$ ; any separating a.f.p. dense sequence in a compact space X with  $|X| \geq 2$  is properly covering.

**Proof.** For each  $x \in X$  we have  $(U_i(x))_1^{r_x}$ ,  $\mathcal{T}_X \ni V_x \ni x$  and  $h_x \in P^{(1)}$  s.t.  $(\forall c \in \nu^{-1}V_x)(\nu(c) \in U_{h_x(c)}(x))$ . By compactness there exist s and  $(x_i)_{i < s} \subseteq X$  with  $X = \cup_{i < s} V_{x_i}$ . We can take a formal disjoint union of  $(U_j(x_i))_{j \in [r_{x_i}]}$  (i < s), say  $(U_i)_{i \in [r]}$  where  $r := \sum_{i < s} r_{x_i}$ , and by adding appropriate constants to each  $h_{x_i}$  (i < s) we get  $h \in P^{(1)}$  with  $(\forall c \in \text{dom } \nu)(\nu(c) \in U_{h(c)})$ .

Proper covering does not imply separation at finite precision.

- **Proposition 12.** 1. Suppose  $\nu_0, \ldots, \nu_n \in \mathcal{I}_d$ ,  $(\forall i < n)(\nu_{i+1} = (\nu_i)_{x_i,y_i}^{A_i})$ ,  $\nu_0$  total & properly covering, and  $\lambda \in \mathcal{I}_d^{rand}$ . Then  $\lambda \not\leq_a \nu_n$ .
  - 2. Suppose  $|X| \geq 2$ . For any  $\nu \in \mathcal{I}_d^{rand}$  there exists  $\lambda \in X^{\mathbb{N}} \setminus \mathcal{I}_d^{rand}$  with  $\nu \leq \lambda$ .
  - 3.  $\mathcal{I}_d^{rand}$  is closed under  $\oplus$ .

**Proof.** (1): Suppose  $\lambda \leq_a \nu_n$  via  $f \in P^{(1)}$  at precision  $\epsilon$ , where  $\epsilon$  suff. small that  $N_{\epsilon}(\{x_i,y_i\})$  and  $N_{\epsilon}(U_j)$  nondense for each i < n and each  $j \in [r]$ , for some open cover  $(U_j)_1^r$ . We have dom  $\nu_n = \mathbb{N} = \dot{\cup}_{i < n} A_i \dot{\cup} 2^n \mathbb{N}$  where each  $A_i$  is infinite c.e. with  $\nu_n(A_i) = \{x_i, y_i\} \subseteq \operatorname{im} \nu_{i+1}$  (i < n). Since  $\lambda$  total,  $f \in R^{(1)}$  with some  $A_i \cap \operatorname{im} f$  (i < n) or  $2^n \mathbb{N} \cap \operatorname{im} f$  infinite (by pigeonhole principle). Each of these sets is c.e. If  $A_i \cap \operatorname{im} f$  infinite, its  $\nu_n$ -image  $(\subseteq \{x_i, y_i\})$  lies within  $\epsilon$  of  $\lambda(f^{-1}A_i)$  which is dense, contradicting choice of  $\epsilon$ . If  $2^n \mathbb{N} \cap \operatorname{im} f$  infinite, let h be as in definition of  $\nu_0$  for cover  $(U_i)_1^r$ . Then

each  $h^{-1}\{i\}$  is c.e., so  $f^{-1}(2^n(h^{-1}\{i\}))$  is c.e., and at least one such set is infinite, so has  $\lambda$ image dense &  $\subseteq N_{\epsilon}(U_i)$  (contradicting choice of  $(U_i)_1^r, h$ ).

(2): For any  $\nu \in \mathcal{I}_{\mathsf{d}}^{\mathsf{rand}} \not\ni \lambda$  we have  $\nu \oplus \lambda \not\in \mathcal{I}_{\mathsf{d}}^{\mathsf{rand}}$ . Such  $\lambda$  can always be found e.g. as a nondense total sequence (extend  $\oplus$  definition). (3): Let  $\lambda_0, \lambda_1 \in \mathcal{I}_{\mathsf{d}}^{\mathsf{rand}}$ . For any denseness test (A, a, j) we have some i < 2 such that  $A \cap (2\mathbb{N} + i)$  is infinite. Clearly  $\lambda_i \leq \lambda_0 \oplus \lambda_1$  via (injective total recursive)  $h: k \mapsto 2k + i$ , and  $B:=h^{-1}(A \cap (2\mathbb{N} + i))$  is infinite c.e. We know  $\lambda_i$  passes the denseness test (B, a, j), say  $k \in B \cap \lambda_i^{-1} \alpha \langle a, j \rangle$ , so  $(\lambda_0 \oplus \lambda_1)(h(k)) = \lambda_i(k)$  shows  $\lambda_0 \oplus \lambda_1$  passes (A, a, j).

Remains to establish  $\mathcal{I}_{d}^{rand} \neq \emptyset$ . For topological space Y, a continuous surjection  $T: Y \rightarrow Y$  is one-sided topologically mixing if

$$(\forall U, V \in \mathcal{T}_Y \setminus \{\emptyset\}) (\exists N \in \mathbb{N}) (\forall n) (n \geq N \implies T^n(U) \cap V \neq \emptyset).$$

### One checks

**Lemma 13.** For any separable metrizable X and  $Y = X^{\mathbb{N}}$ , left shift  $\sigma : Y \to Y$  is one-sided topologically mixing ( w.r.t. product topology).

More generally, consider a complete effective metric space  $(Y,d,\nu)$  with ideal ball numbering  $\alpha$  and continuous one-sided top. mixing  $T:Y\to Y$ . By definition, each  $\cup_{m\in A}T^{-m}V$  dense  $(V\in\mathcal{T}_Y\setminus\{\emptyset\},\ A \text{ infinite})$ , so in particular  $R_A:=\cap_{a\in\text{dom }\alpha}\cup_{m\in A}T^{-m}\alpha(a)$  and  $R:=\cap\{R_A\mid A \text{ infinite c.e.}\}$  dense  $\mathcal{G}_\delta$  in Y (by Baire category theorem). We now apply to  $Y=X^\mathbb{N}$ . Recall a formal inclusion  $\square$  of total basis numberings  $\alpha$ ,  $\beta$  of space X has the weak basis property if  $(\forall b)(\forall x\in X)(\exists a)(x\in\beta(b)\Longrightarrow x\in\alpha(a)\land a\sqsubseteq b)$ .

**Proposition 14.** Let  $(X,d,\nu_0)$  be a complete bounded effective metric space with  $\nu_0$  total, and equip  $X^{\mathbb{N}}$  with the (bounded) product metric  $\widehat{d}(\xi,\eta) := \sum_{i \in \mathbb{N}} 2^{-i-1} d(\xi_i,\eta_i)$  and dense sequence  $\gamma: \mathbb{N} \to X^{\mathbb{N}}$  defined by  $\gamma(\langle w \rangle)(i) := \begin{cases} \nu_0(w_i), & \text{if } i < |w|, \\ \nu_0(i), & \text{if } i \geq |w| \end{cases}$ . Then basis numberings defined by

$$lpha_{X^{\mathbb{N}}}: \mathbb{N} \to \mathcal{T}_{X^{\mathbb{N}}}, \langle a, r \rangle \mapsto B_{\widehat{d}}(\gamma(a); \nu_{\mathbb{Q}^{+}}(r)),$$

$$\beta: \mathbb{N} \to \mathcal{T}_{X^{\mathbb{N}}}, \langle a, r, N \rangle \mapsto \prod_{i < N} B_{d}(\gamma(a)_{i}; \nu_{\mathbb{Q}^{+}}(r)) \times X^{\mathbb{N}}$$

are effectively equivalent, and any  $\xi \in X^{\mathbb{N}}$  has  $\left((\sigma^l \xi)_{l \in \mathbb{N}} \in \mathcal{I}_d^{rand}|_{X^{\mathbb{N}}} \iff \xi \in R|_{X^{\mathbb{N}},\sigma} \implies \xi \in \mathcal{I}_d^{rand}|_{X}\right);$  in particular  $\mathcal{I}_d^{rand}|_{X^{\mathbb{N}}}$  is residual in  $X^{\mathbb{N}}$ .

**Proof.** It is convenient to use the following binary relations on  $\mathbb{N}$ :

$$\langle a,r \rangle \sqsubseteq \langle b,q,N \rangle : \iff \max_{i < N} d(\gamma(a)_i,\gamma(b)_i) + 2^{i+1} \nu_{\mathbb{Q}^+}(r) < \nu_{\mathbb{Q}^+}(q),$$
 $\langle a,r,N \rangle \sqsubseteq' \langle b,q \rangle : \iff \sum_{i < N} 2^{-i-1} (\nu_{\mathbb{Q}^+}(r) + d(\gamma(a)_i,\gamma(b)_i))$ 
 $+ C_0.2^{-N} < \nu_{\mathbb{Q}^+}(q);$ 

here  $\mathbb{Q}^+ \ni C_0 \ge \operatorname{diam}(X,d)$  is a fixed bound. We first check  $\sqsubseteq$ ,  $\sqsubseteq'$  are c.e. formal inclusions (of  $\alpha_{X^{\mathbb{N}}}$  in  $\beta$ , resp. of  $\beta$  in  $\alpha_{X^{\mathbb{N}}}$ ), each with the weak basis property.

Next denote  $B_{\langle a,r\rangle}:=\{\langle b,q\rangle\,\big|\,\langle b,q\rangle\, \sqsubset \langle\langle a\rangle,r,1\rangle\}$   $(a,r\in\mathbb{N})$ . Clearly  $B_{\langle a,r\rangle}$  is nonempty and c.e. uniformly in  $a,r\in\mathbb{N}$ , and for any  $A\subseteq\mathbb{N}$  the condition  $(\forall a,r)\, \big(\xi\in\cup_{l\in A}\sigma^{-l}(\alpha\langle a,r\rangle\times X^{\mathbb{N}})\big)$  is equivalent to  $(\forall a,r)(\exists\langle b,q\rangle\in B_{\langle a,r\rangle})\, \big(\xi\in\cup_{l\in A}\sigma^{-l}\alpha_{X^{\mathbb{N}}}\langle b,q\rangle\big)$ . In particular, this is implied by  $(\forall b,q)\, \big(\xi\in\cup_{l\in A}\sigma^{-l}\alpha_{X^{\mathbb{N}}}\langle b,q\rangle\big)$ , or equivalently  $(\forall b,q)(\exists k\in A)\, \big((\sigma^{l+k}\xi)_{l\in\mathbb{N}}\in\alpha_{X^{\mathbb{N}}}\langle b,q\rangle\times (X^{\mathbb{N}})^{\mathbb{N}}\big)$ . Quantifying over infinite c.e. A, thus  $(\xi\in R|_{X^{\mathbb{N}},\sigma}\Longleftrightarrow)$   $(\sigma^l\xi)_{l\in\mathbb{N}}\in\mathcal{I}_{\mathsf{d}}^{\mathsf{rand}}|_{X^{\mathbb{N}}}$  implies  $\xi\in\mathcal{I}_{\mathsf{d}}^{\mathsf{rand}}|_{X^{\mathbb{N}}}$ 

**Example:**  $X = \{0,1\}$  corresponds (easily checked) to  $\mathcal{I}_{\mathsf{d}}^{\mathsf{rand}}|^X = \{\chi_B \, \big| \, B \subseteq \mathbb{N} \, \text{ bi-immune} \}$ , i.e. those sets B for which neither B nor  $\mathbb{N} \setminus B$  contains an infinite c.e. set.

# Comparison with other tests

For any  $A\subseteq \mathbb{N}$ ,  $N\in \Pi^0_1(X)$  denote

$$F_{A,N} := \prod_{i \in \mathbb{N}} W_i \text{ where } W_i = \begin{cases} X, & \text{if } i \not \in A, \\ N, & \text{if } i \in A \end{cases}.$$

 $F_{A,X\setminus lpha\langle a,j
angle}=\{\xi\,ig|\,\xi \ ext{fails}\ (A,a,j)\} \ ext{for denseness test}\ (A,a,j),$   $X^{\mathbb{N}}\setminus \mathcal{I}^{\mathsf{rand}}_{\mathsf{d}}=\cup \{F_{A,N}\,ig|\,A \ ext{infinite c.e.,}\ N=X\setminus lpha\langle a,j
angle, a,j\in\mathbb{N}\,\}.$ 

??For X effectively compact,

$$f: \{0,1\}^{\mathbb{N}} \times \mathcal{K}_{>}(X) \to \mathcal{K}_{>}(X^{\mathbb{N}}), (\chi_{A},K) \mapsto F_{A,K}$$

computable. More generally, suppose  $|X| \geq 2$ . If  $\mu_0$  is a measure positive on nonempty open sets, for any test (A,a,j) we have  $\mu_0(X\backslash B_0) < 1$  for  $B_0 := \alpha \langle a,j \rangle$ , so  $F_{A,X\backslash B_0}$  is a closed  $\lambda$ -nullset. Here  $\lambda$  is the product measure on  $X^\mathbb{N}$  corresponding to measure  $\mu_0$  on X; also one can show  $\lambda$  a computable probability measure if  $\mu$  is and X bounded complete.

### Effectiveness and approximability

When choosing naming system  $\gamma$  for a given space Y, several types of requirements [?,  $\S 2.7$ ]:

- (a) require  $f_i :\subseteq X_i \to Y$  computable  $(i \in I)$
- (b) require  $f_i :\subseteq Y \to Z_i$  computable  $(i \in I)$
- (c) require operations  $f_i:\subseteq Y^{n_i}\to Y$  computable  $(i\in I)$

For example, types (a)+(c) and (b) resp. addressed by **Proposition 15.** Fix represented sets  $(X_i, \rho_i)$ , maps  $f_i :\subseteq X_i \to Y$ ,  $g_i :\subseteq Y \to Y$   $(i \in \mathbb{N})$  & suppose

 $Y=\cup_{i\in\mathbb{N}\,,w\in\mathbb{N}^*}\widehat{g}_w(\mathsf{im}\,f_i\cap\mathsf{dom}\,\widehat{g}_w)$  where  $\widehat{g}_w:=g_{w_{|w|-1}}\circ\ldots\circ g_{w_0}$ .

 $\delta :\subseteq \mathbb{N}^{\mathbb{N}} \to Y$  defined by  $\operatorname{dom} \delta = \dot{\cup}_{i \in \mathbb{N}, w \in \mathbb{N}^*} i.\langle w \rangle.(\rho_i^{-1} f_i^{-1} \operatorname{dom} \widehat{g}_w)$  and  $\delta(i.\langle w \rangle.p) := (\widehat{g}_w \circ f_i \circ \rho_i)(p)$  has  $\delta \in \mathcal{R}_Y$ .  $f_i$  is  $(\rho_i, \delta)$ -computable and  $g_i$  is  $(\delta, \delta)$ -computable  $(i \in \mathbb{N})$ .

**Proposition 16.** Fix represented sets  $(Y_i, \delta_i)$ ,  $f_i :\subseteq X \to Y_i$   $(i \in \mathbb{N})$  where  $|X| \leq 2^{\aleph_0}$ . Then there exists  $\rho \in \mathcal{R}_X$  such that each  $f_i$  is strongly  $(\rho, \delta_i)$ -computable  $(i \in \mathbb{N})$ .

More sophisticated results for type (c), see [?, Thm 2.7.15], [?],

For requirements of type (a)+(b) or (b)+(c), in general there exist counterexamples:

- **Proposition 17.** 1. there exist sets X,Y,Z, maps  $f:\subseteq X \to Y$ ,  $g:\subseteq Y \to Z$  and total numberings  $\nu,\mu$  of X,Z respectively, such that  $\neg(\exists \lambda:\subseteq \mathbb{N} \to Y)(f\circ \nu \leq \lambda \wedge g\circ \lambda \leq \mu)$ .
  - 2. there exist sets Y, Z, maps  $f :\subseteq Y \to Y$ ,  $g :\subseteq Y \to Z$  and total numbering  $\mu$  of Z such that  $\neg(\exists \lambda \in \mathsf{TN}(Y))(f \circ \lambda \leq \lambda \land g \circ \lambda \leq \mu)$ .
  - 3. there exist separable metric spaces Y, Z, maps  $f : \subseteq Y \to Y$ ,  $g : \subseteq Y \to Z$  and  $\mu \in \mathcal{I}_d^Z \cap Z^{\mathbb{N}}$  such that  $\neg (\exists \lambda \in \mathcal{I}_d^Y) (f \circ \lambda \leq \lambda \wedge g \circ \lambda \leq \mu)$ .
- **Proof.** (1): X = Y = Z denumerable,  $f = g = \mathrm{id}_X$ ,  $\nu$  total injective numbering of X,  $\mu$  a total numbering s.t.  $\nu \not \leq \mu$  (e.g. an incomparable total injective numbering, of which there are  $2^{\aleph_0}$ ; see [?]).
- (2): Y = Z denumerable,  $\mu$  a total injective numbering,  $h: \mathbb{N} \to \mathbb{N}$  a nonrecursive bijection with  $N_h: \mathbb{N} \to \overline{\mathbb{N}}$ ,  $k \mapsto \#\operatorname{Fix}(h^k)$  not lower semicomputable,  $f = \mu \circ h \circ \mu^{-1}$ ,  $g = \operatorname{id}_Y$ . If  $\lambda \in \operatorname{TN}(Y)$  and  $n \in P^{(1)}(\lambda, \mu)$ -realises  $g, =_Y$  is  $[\lambda, \lambda]$ -decidable via  $\langle a, b \rangle \mapsto \delta_{n(a), n(b)}$  (since  $g, \mu$  injective). So  $N_{f,M}: \mathbb{N} \to \mathbb{N}$ ,  $k \mapsto \#\operatorname{Fix}(f^k|_{\lambda[0,M)})$   $(M \ge 1)$  are uniformly effective provided also f is  $(\lambda, \lambda)$ -effective: consider

$$a_n := \mu a < M\left((f^k \circ \lambda)(a) = \lambda(a) \notin \{\lambda(a_i) \mid i < n\}\right)$$

with convention  $\mu a < M$  (False) = M, inductively for  $n \le M$  and  $N := \mu n$  ( $a_n = M$ ) ( $\le M$ ). Then  $\lambda|_{\{a_i|i < N\}}$  is injective with image Fix( $f^k|_{\lambda[0,M)}$ ), so  $N_{f,M}(k) = N$ .

Further,  $(f^k \circ \lambda)(a) = \lambda(a)$  iff  $(h^k \circ \mu^{-1} \circ \lambda)(a) = (\mu^{-1} \circ \lambda)(a)$ , so  $\operatorname{Fix}(h^k|_{(\mu^{-1}\circ\lambda)[0,M)}) = \mu^{-1}\operatorname{Fix}(f^k|_{\lambda[0,M)})$ . In unions over increasing M,  $\operatorname{Fix}(h^k) = \mu^{-1}\operatorname{Fix}(f^k)$  and  $N_h(k) = \sup_M N_{f,M}(k)$ . This contradicts choice of h.

(3):  $Y = Z = \mathbb{R}$ ,  $f = \mathrm{id}_{\mathbb{R}} + \alpha$  ( $\alpha \in \mathbb{R} \setminus \mathbb{Q}$ ),  $g = \mathrm{id}_{\mathbb{R}}$ ,  $\mu = \nu_{\mathbb{Q}}$ . The only solution for  $\lambda :\subseteq \mathbb{N} \to Y$  is then the empty partial sequence.

Next, consider (a),(b),(c) when maps are approximable rather than computable. Here  $f:\subseteq X \to Y$  approximable w.r.t.  $\nu \in \mathcal{I}_d^X$ ,  $\lambda \in \mathcal{I}_d^Y$  if  $f \circ \nu \leq_a \lambda$  (extending definition of  $\leq_a$ ).

If X uniformly discrete, (1), (2) transfer since  $\nu \leq \lambda$  iff  $\nu \leq_a \lambda$  for any  $\nu, \lambda \in \mathcal{I}_d^X = \mathsf{TN}(X)$ . If X compact, instead take  $\nu_0 \in \mathcal{I}_d^X$  compact and pick  $\nu, \mu \leq_a$ -incomparable. Then  $\neg(\exists \lambda : \subset \mathbb{N} \to Y)(f \circ \nu <_a \lambda \land g \circ \lambda <_a \mu)$ .

**Lemma 18.** If X compact,  $\nu \in \mathcal{I}_d$  compact and  $\pi : X \to Y$  a continuous surjection then  $\pi \circ \nu$  is compact.

**Proof.** Uses Lebesgue numbers and uniform continuity of  $\pi$ ; see [?].

**Corollary 19.** Let Y, Z be compact metric spaces,  $f: Y \to Y$ ,  $g: Y \to Z$  (total) continuous surjective and  $\mu \in \mathcal{I}_d^Z$ . For any compact  $\lambda \in \mathcal{I}_d^Y$  we have  $f \circ \lambda \leq_a \lambda \land g \circ \lambda \leq_a \mu$ .

**Proof.** Follows from Lemma 18 and analogue of Lemma 1.

Given existence of compact  $\lambda \in \mathcal{I}_{d}^{Y}$ , this is a positive result for requirements of type (b)+(c).

# Requirements of type (a),(b)

For  $\nu \in \mathcal{I}_X$  we denote the cylindrification by  $\tilde{\nu}$ : dom  $\tilde{\nu} = \langle \operatorname{dom} \nu, \mathbb{N} \rangle$ ,  $\tilde{\nu} \langle i, j \rangle := \nu(i)$ .

**Proposition 20.** Fix separable metric spaces  $X, Y_i$ , maps  $\Psi_i : X \to Y_i$  and  $\lambda_i \in \mathcal{I}_d^{Y_i}$  (i < n). There exists  $\nu \in \mathcal{I}_d^X$  s.t.  $(\forall i < n) \Psi_i \circ \nu \leq_a \lambda_i$ .

**Proof.** Fix  $\nu_0 \in \mathcal{I}_d \cap X^{\mathbb{N}}$  and (for  $i < n, k \in \mathbb{N}$ )

$$a_k^{(i)} := \mu n (n \in \operatorname{dom} \tilde{\lambda}_i \wedge \neg (\exists j < k) (n = a_j) \wedge d((\Psi_i \circ \nu_0)(k), \tilde{\lambda}_i(n)) < 2^{-k}).$$

We denote  $\nu\langle a_k^{(0)},\ldots,a_k^{(n-1)}\rangle:=\nu_0(k)$ , with  $\nu$  undefined elsewhere. For any  $\epsilon>0$ ,  $\max_{i< n}d((\Psi_i\circ\nu)\langle a_k^{(0)},\ldots,a_k^{(n-1)}\rangle,\tilde{\lambda}_i(a_k^{(i)}))<\epsilon$  fails only for finitely many k (among those with  $\epsilon\leq 2^{-k}$ ). We also have  $\operatorname{im}\nu=\operatorname{im}\nu_0$  and  $\tilde{\lambda}_i\equiv\lambda_i$  for each i< n.

**Proposition 21.** Let X,Y be separable metric spaces,  $\Psi: X \to Y$ ,  $\nu \in \mathcal{I}_d^X$ ,  $\lambda \in \mathcal{I}_d^Y$  as appropriate. Then:

- 1. for any  $\nu$  there exists  $\lambda$  with  $\Psi \circ \nu \leq \lambda$ ; if  $\Psi$  continuous onto Y, there exists  $\lambda$  with  $\Psi \circ \nu \equiv \lambda$ ,
- 2. if  $\Psi$  open, for any  $\lambda$  there exists  $\nu$  with  $\Psi \circ \nu \leq \lambda$ ,
- 3. for any  $\nu$  there exists  $\lambda$  with  $\Psi \circ \nu \leq_a \lambda$ ,
- 4. for any  $\lambda$  there exists  $\nu$  with  $\Psi \circ \nu \leq_a \lambda$ .

**Proof.** (2): Each  $\Psi^{-1}\{\lambda(a)\}$  separable  $(a \in \text{dom }\lambda)$  so if nonempty fix a dense subsequence  $(x_i^{(a)})_{i \in \mathbb{N}}$  and let  $\nu\langle a,i\rangle := x_i^{(a)}$   $(i \in \mathbb{N})$   $(\nu$  undefined elsewhere). Since  $\Psi$  open,  $\Psi^{-1}$  im  $\lambda$  is dense, hence so is im  $\nu$ .

For (4), if  $\Psi$  onto we can construct  $\nu \in \mathcal{I}_{d}^{X}$  in a different way. First, pick  $\nu_{0} \in \mathcal{I}_{d}^{X} \cap X^{\mathbb{N}}$  and  $\lambda \in \mathcal{I}_{d}^{Y}$ . Let  $\operatorname{dom} \nu = \{a_{k} \mid k \in \mathbb{N}\}, \ \nu(a_{k}) = \nu_{0}(k)$  for  $a_{k} := \mu n \ (n \in \operatorname{dom} \tilde{\lambda} \wedge \neg (\exists j < k) (n = a_{j}) \wedge \tilde{\lambda}(n) \in U_{k}) \quad (k \in \mathbb{N})$ 

where  $(U_k)_{k\in\mathbb{N}}\subseteq\mathcal{T}_Y\setminus\{\emptyset\}$  has the property  $(\forall x\in \operatorname{im}\lambda)(\forall N)(\exists k\geq N)(U_k\ni x)$ . In place of  $\operatorname{dom}\nu\subseteq \operatorname{dom}\tilde{\lambda}$  we now have equality: for  $c_l$  the  $l^{\operatorname{th}}$  element of  $\operatorname{dom}\tilde{\lambda}$  in ascending order, and  $N_l:=\mu k\left(U_k\ni\tilde{\lambda}(c_l)\wedge(\forall l'< l)(k>N_{l'})\right)$ , we show  $c_l\in\{a_j\,|\,j\leq N_l\}$   $(l\in\mathbb{N})$ . First,  $N_0=\mu k\left(U_k\ni\tilde{\lambda}(c_0)\right)$ , so  $\neg(\exists j< N_0)(c_0=a_j)$  (as  $\tilde{\lambda}(a_j)\in U_j$ ) and then  $a_{N_0}=c_0$ . Secondly,  $a_{N_{l+1}}=\min(\tilde{\lambda}^{-1}U_{N_{l+1}})\setminus\{a_j\,|\,j< N_{l+1}\}$  and  $U_{N_{l+1}}\ni\tilde{\lambda}(c_{l+1})$  so if  $c_{l+1}\not\in\{a_j\,|\,j\leq N_{l+1}\}$  then  $a_{N_{l+1}}< c_{l+1}$ , implying  $(\exists k\leq l)a_{N_{l+1}}=c_k\in\{a_j\,|\,j\leq N_k\}$ . But then  $N_{l+1}>N_l\geq N_k$  contradicts injectivity of  $(a_i)_{i\in\mathbb{N}}$ .

To ensure also  $\Psi \circ \nu \equiv_a \tilde{\lambda}$ , we assume compactness. We also note two improvements on Proposition 21(2) with similar conditions.

**Proposition 22.** Suppose X, Y are separable metric spaces,  $\lambda \in \mathcal{I}_d^Y$ ,  $\Psi : X \to Y$  any map.

1. If  $\Psi$  onto and Y compact, there exists  $\nu \in \mathcal{I}_d^X$  such that  $\Psi \circ \nu \equiv_a \lambda$ ,

- 2. If  $\overline{\Psi^{-1}}$  im  $\lambda = X$  and each  $\Psi^{-1}\{y\}$  is compact  $(y \in \operatorname{im} \lambda)$ , there exists  $\nu \in \mathcal{I}_d^X$  such that  $\Psi \circ \nu = \tilde{\lambda}$ .
- 3. If  $\overline{\Psi^{-1}}$  im  $\lambda = X$  and X compact, there exists compact  $\nu \in \mathcal{I}_d^X$  such that  $\Psi \circ \nu \leq \lambda$ .

**Proof.** (1): In view of the above, it remains to construct  $(U_k)_k$  appropriately. As Y compact, let  $\nu_0$  be such that  $\nu_0|_{2\mathbb{N}+1}$  is dense and  $2(\sum_{j< l} n_j + i) \in \nu_0^{-1} \Psi^{-1} \{y_i^l\}$   $(l \in \mathbb{N}, i < n_l)$  where  $\{y_i^l \mid i < n_l\}$  is a  $2^{-l}$ -spanning set for Y. Then let  $U_k := B((\Psi \circ \nu_0)(k); 2^{-b_k})$  where  $b_k = l$  if  $2\sum_{j< l} n_j \leq k < 2\sum_{j\leq l} n_j$ . This ensures  $\lim_{k \to \infty} b_k = \infty$  and so  $\Psi \circ \nu \equiv_a \tilde{\lambda}$ .

(2): For each  $a \in \text{dom } \lambda$ ,  $k \in \mathbb{N}$  there is a finite  $2^{-k}$ -spanning subset of  $\Psi^{-1}\{\lambda(a)\}$ , say  $(z_i(a,k))_{i < m(a,k)}$ . Let  $\nu\langle a, \sum_{j < k} m(a,j) + i \rangle := z_i(a,k)$  for each i < m(a,k) (also  $\nu\langle a,m \rangle \uparrow$  if  $a \not\in \text{dom } \lambda$ ). Plainly  $(\Psi \circ \nu)\langle a,m \rangle = \lambda(a) = \tilde{\lambda}\langle a,m \rangle$  whenever  $a \in \text{dom } \lambda \wedge m \in \mathbb{N}$ , and both sides undefined for

 $a \not\in \text{dom }\lambda$ . Given  $x \in X$ ,  $k \in \mathbb{N}$  there exist  $a \in \text{dom }\lambda$ ,  $y \in \Psi^{-1}\{\lambda(a)\} \cap B(x; 2^{-k-1})$  and i < m(a,k+1) s.t.  $z_i(a,k+1) \in B(y; 2^{-k-1}) \cap \text{im }\nu$ . Then  $d_{\text{im }\nu}(x) < 2^{-k-1} + 2^{-k-1}$ , and k was arbitrary.

We describe a general extension of (3) (of type (b)+(c)):

**Proposition 23.** Let X,Y be separable metric spaces, X compact,  $\lambda \in \mathcal{I}_d^Y$ ,  $g :\subseteq X \to Y$  any map with  $\overline{g^{-1} \operatorname{im} \lambda} = X$ ,  $\mathcal{F}$  a denumerable family of (total) maps  $X \to X$ . Also take  $(\theta_k)_{k \in \mathbb{N}} \subseteq (0, \infty)$  with  $\inf_k \theta_k = 0$ ,  $E_l \subseteq g^{-1} \operatorname{im} \lambda$  a finite  $\theta_l$ -spanning set  $(l \in \mathbb{N})$ . Suppose  $\sim_l (l \in \mathbb{N})$  are equivalence relations on  $\mathcal{F}$  s.t.  $\mathcal{F}/\sim_l$  finite,

$$(\forall y \in E_l)(\forall C \in \mathcal{F}/\sim_l)(\forall i \leq l)(\exists z \in E_i)\{fy \mid f \in C\} \subseteq B(z; \theta_i),$$
(2)

and which are increasingly fine & satisfy separating condition  $(\forall f, f' \in \mathcal{F})(\exists l)(f \neq f' \implies f \not\sim_l f')$ . Then there exists compact  $\nu \in \mathcal{I}_d^X$  with  $g \circ \nu \leq \lambda$  and  $(\forall f \in \mathcal{F})(f \circ \nu \leq_a \nu)$ .

**Proof.** Write  $E_k = \{y_1^k, \dots, y_{r_k}^k\} \subseteq g^{-1} \text{ im } \lambda$  (a  $\theta_k$ -spanning set for X,  $k \in \mathbb{N}$ ). Based on  $\sim_l$   $(l \in \mathbb{N})$  and an injective total numbering  $e \mapsto f_e$  of  $\mathcal{F}$ , we define  $\nu$  with  $\text{im } \nu = \cup_k E_k$ ; surjectivity will follow from (2), density from  $\text{inf}_k \theta_k = 0$ . Separating condition ensures the partitions  $\mathcal{P}_l := \mathcal{F}/\sim_l (l \in \mathbb{N})$  (with  $(\forall l)\mathcal{P}_{l+1} \geq \mathcal{P}_l$ ) are generating:  $|\cap_l A_l| \leq 1$  for any  $A_l \in \mathcal{P}_l$   $(l \in \mathbb{N})$ . If each  $\mathcal{P}_l$  is endowed with a numbering  $\nu_l :\subseteq \mathbb{N} \to \mathcal{P}_l$ , this leads to a representation  $\rho :\subseteq \mathbb{N}^{\mathbb{N}} \to \mathcal{F}$ :  $\text{dom } \rho = \{p \in \mathbb{N}^{\mathbb{N}} \mid (\forall l)(p_l \in \text{dom } \nu_l) \wedge \cap_l \nu_l(p_l) \neq \emptyset\}$ ,  $\rho(p) \in \cap_l \nu_l(p_l)$ .

Will return to numbering  $e\mapsto f_e$  and representation  $\rho$  of  $\mathcal{F}$ . Requirement for  $f\in\mathcal{F}$  to be  $(\nu,\nu)$ -approximable equivalent to  $f\circ\nu$  compact — (roughly speaking) suggests to store approximations to each f in each  $\nu$ -name a. Form of  $\nu$  will be

$$a = \langle l; j_0, \dots, j_l; j_0^{(0)}, \dots, j_l^{(0)}; \dots; j_0^{(m)}, \dots, j_l^{(m)}; \hat{b} \rangle \in \text{dom } \nu,$$
  
 $\nu(a) = y_{j_l}^l;$ 

 $(\forall i < l)d(y_{j_i}^i, y_{j_i}^l) < \theta_i$ ,  $\hat{b} = b(l, j_l)$  for a fixed choice function  $b :\subseteq \mathbb{N}^2 \to \mathbb{N}$  with  $\operatorname{im} b \subseteq \operatorname{dom} \lambda$ ,  $\operatorname{dom} b =$ 

 $\{(k,j) \mid 1 \leq j \leq r_k\}$  and  $g(y_j^k) = (\lambda \circ b)(k,j)$  for all  $(k,j) \in \text{dom } b$ , plus further conditions on m and  $j_i^{(q)}$   $(q \leq m, i \leq l)$ . Namely, assuming dom  $\nu_l$  finite, want  $m \geq \max(\text{dom } \nu_l)$  and

$$d(fy_{j_l}^l,y_{j_i^{(p_l)}}^i)< heta_i$$
 whenever  $i\leq l\in\mathbb{N}\,,f\in\mathcal{F},p\in
ho^{-1}\{f\}.$ 

One checks existence of suitable  $j_i^{(q)}$   $(q \in \text{dom } \nu_l, i \leq l)$  for given  $y = y_{j_l}^l \in E_l$  follows from (2).

For simplicity let  $\operatorname{dom} \nu_l = [0, m_l]$  and  $m_l := |\mathcal{P}_l| - 1$   $(l \in \mathbb{N})$ , with  $m = m_l$  determined by l. To complete the construction, there is some freedom in choice of  $\nu_l$ . We will ensure each  $f \in \mathcal{F}$  has a computable  $\rho$ -name; the (l+1)-block  $j_0^{(q)}, \ldots, j_l^{(q)}$  relevant to f can be picked out of  $\nu$ -name a based on this. More rigourously, fix choice function  $c'' : \subseteq \mathcal{F} \times \mathbb{N}^3 \to \mathbb{N}$  s.t.

$$d(fy_j^l, y_{c''(f,l,k,j)}^k) < \theta_k$$
 for all  $(f,l,k,j) \in \text{dom } c'' := \{(f,l,k,j) \mid f \in \mathcal{F} \land l < k \land 1 \leq j \leq r_l\}.$ 

For any fixed k and  $f \in \mathcal{F}$ , from  $p \in \rho^{-1}\{f\}$  and  $a \in \text{dom } \nu$  we will be able to compute

$$u := \begin{cases} j_k^{(p_l)}, & \text{if } k \leq l, \\ c''(f, l, k, j_l), & \text{if } l < k, \end{cases}$$

and observe  $d(fy_{j_l}^l, y_u^k) < \theta_k$ ; this computation is possible since  $(\{f\} \times \mathbb{N} \times \{k\} \times \mathbb{N}) \cap \text{dom } c''$  is finite.

To define  $\nu_l$   $(l \in \mathbb{N})$ , proceed in stages; at the end of stage i we will have each  $\nu_l$  defined on an interval  $[0,m_{l,i}]$  with

$$\nu_l[0, m_{l,i}] = \{ [f_j]_{\sim_l} | j \le i \}$$
 for all  $l \in \mathbb{N}$ . (3)

In stage 0 we define  $m_{l,0} := 0$  and declare  $0 \in \nu_l^{-1}\{[f_0]_{\sim_l}\}$  for all l; also let  $l_0 := 0$  and  $p^{(0)} := 0^{\omega} \ (\in \rho^{-1}\{f_0\})$ .

At stage i+1 let  $l_{i+1} := \inf\{l \in \mathbb{N} \mid (\forall j \leq i)(f_{i+1} \not\sim_l f_j)\}$   $(i \in \mathbb{N})$ . For each  $l < l_{i+1}$  we have  $f_{i+1} \in \cup_{j \leq i} [f_j]_{\sim_l}$ , so  $(\exists p_l^{(i+1)} \leq m_{l,i})(\nu_l(p_l^{(i+1)}) = [f_{i+1}]_{\sim_l})$ , and we leave  $\nu_l$  unmodified,  $m_{l,i+1} := m_{l,i}$ . For  $l \geq l_{i+1}$  we have  $[f_{i+1}]_{\sim_l} \not\in \{[f_j]_{\sim_l} \mid j \leq i\} = \nu_l[0, m_{l,i}]$ , so let  $p_l^{(i+1)} := m_{l,i+1} := m_{l,i} + 1$  and  $\nu_l(m_{l,i+1}) := [f_{i+1}]_{\sim_l}$ .

In either case  $(3)|_{i+1}$  holds (by inspection), computability of  $(m_{l,i+1})_l \in \mathbb{N}^{\mathbb{N}}$  holds by induction, and the computability of  $p^{(i+1)} \in \rho^{-1}\{f_{i+1}\}$  will follow once we show  $l_{i+1} < \infty$ . Suppose

that  $l_{i+1}=\infty$ ; then  $(\forall l)(\exists j\leq i)(f_{i+1}\sim_l f_j)$ . By the pigeonhole principle and the fact  $(\mathcal{P}_l)_l$  are increasingly fine, we get  $(\exists j\leq i)(\forall l)(f_{i+1}\sim_l f_j)$ , so  $f_{i+1}=f_j$  since  $(\mathcal{P}_l)_l$  is generating. This contradicts injectivity of  $e\mapsto f_e$ .

Finally, since  $\{f_e \mid e \in \mathbb{N}\} = \mathcal{F}$ ,  $(\forall i)(3)|_i$  implies  $\nu_l(\cup_i[0,m_{l,i}]) = \mathcal{P}_l$  for each l. It is clear by construction  $\nu_l$  injective, so  $\sup_i m_{l,i}$  finite, and  $\sup_i m_{l,i} = m_l = |\mathcal{P}_l| - 1$  as previously assumed.

Clear how to check compactness of  $\nu$  w.r.t.  $(B(y_j^k;\theta_k))_{j=1}^{r_k}$ : use choice function  $c^+:\subseteq \mathbb{N}^3 \to \mathbb{N}$  with  $d(y_j^l,y_{c^+(l,k,j)}^k)<\theta_k$  for all  $(l,k,j)\in \operatorname{dom} c^+:=\{(l,k,j) \mid l< k, 1\leq j\leq r_l\}$ , and from  $a\in \operatorname{dom} \nu$  compute

$$u := (j_k, \text{ if } l \ge k; c^+(l, k, j_l), \text{ if } l < k);$$

this uses finiteness of  $\mathbb{N} \times \{k\} \times \mathbb{N}$ )  $\cap$  dom  $c^+$ .

If  $id_X \in \mathcal{F}$ , this argument (and data  $j_0, \ldots, j_{l-1}$  in names  $a \in dom \nu$ ) can be omitted. Despite these good properties, several reasons to simplify or generalise above proof.

# Examples for $\sim_l$

Consider equivalence relations  $\sim_l (l \in \mathbb{N})$  defined from choice function  $\tilde{c}$  as follows: fix  $\tilde{c}:\subseteq \mathcal{F}\times\mathbb{N}^3\to\mathbb{N}$  such that

$$d(fy_j^l,y_{\widetilde{c}(f,l,i,j)}^i) < heta_i \quad ext{ for all } (f,l,i,j) \in ext{dom } \widetilde{c} := \{(f,l,i,j) \ ig| \ i \leq l, 1 \leq j \leq r_l\},$$

define  $f \sim_l f'$  if  $\tilde{c}(f,\cdot), \tilde{c}(f',\cdot)$  agree on  $([0,l] \times \mathbb{N}^2) \cap \text{dom } \tilde{c}(f,\cdot)$ . Such  $\sim_l$  are increasingly fine, have  $\mathcal{F}/\sim_l$  finite and (2) holds (definition of  $\nu$  can further be 'simplified' by requiring strictly  $j_i^{(q)} = \tilde{c}(f,l,i,j_l)$  for all  $i \leq l$  whenever  $f \in \nu_l(q), q \leq m_l$ ).

If  $f, f' \in \mathcal{F}$  distinct over  $\cup_k E_k$ , pick  $k, x \in E_k$  s.t.  $f(x) \neq f'(x)$ , and w.l.o.g. assume  $(\forall l)(E_l \subseteq E_{l+1})$ . For  $l \geq k$  s.t.  $2\theta_l < d(fx, f'x)$  and  $j \in [r_l]$  s.t.  $y_j^l = x$  we must have  $\tilde{c}(f, l, l, j) \neq \tilde{c}(f', l, l, j)$ , hence  $f \not\sim_l f'$ . So the separating condition may be replaced by

$$(\forall f, f' \in \mathcal{F})(f \neq f' \implies f|_{\bigcup_k E_k} \neq f'|_{\bigcup_k E_k}),$$

and  $\cup_k E_k$  chosen as any dense countable subset of  $g^{-1}$  im  $\lambda$  (while assuming  $(\forall l)(E_l \subseteq E_{l+1})$ ).

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